

Valószínűség-maximalizálás belső közelítéssel

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Outline

Probabilistic optimization problems

Classic and recent solution approaches

A variant of a classic approach: epi-approximation

Approximating the epigraph of the probabilistic function

Computational study

Theoretical justification of heuristic adjustments

Working with gradient estimates

Problem formulation

ξ an n -dimensional random vector, with known distribution.

Let $F(\cdot)$ denote the cumulative distribution function.

$\mathbf{x} \in \mathbb{R}^m$ decision variables

$T\mathbf{x}$ a linear function of the decision variables

Probabilistic function

$$P(T\mathbf{x} \geq \xi) = F(T\mathbf{x})$$

Probability maximization

$$\max F(T\mathbf{x}) \quad \text{subject to} \quad \mathbf{x} \in X.$$

Joint probabilistic constraint

$$\min \mathbf{c}^T \mathbf{x} \quad \text{subject to} \quad \mathbf{x} \in X, \quad F(T\mathbf{x}) \geq p.$$

Equivalent formulation for joint probabilistic constraint:

$$\min \mathbf{c}^T \mathbf{x} \quad \text{subject to} \quad \mathbf{x} \in X, \quad T\mathbf{x} \in \mathcal{L},$$

$$\text{where} \quad \mathcal{L} = \{ \mathbf{z} \mid F(\mathbf{z}) \geq p \}.$$

Early application

strategic planning model for the Hungarian energy sector
Prékopa, Ganczer, Deák, Patyi (1980).

Characterization and convexity statements, Prékopa (1970-1973):

Assume ξ has a continuous distribution with a log-concave density function.

\Rightarrow the cumulative distribution function $F(\cdot)$ is log-concave

\Rightarrow the probabilistic function $\boldsymbol{x} \mapsto F(T\boldsymbol{x})$ is log-concave.

Solution approaches

Feasible direction method in the convex level set \mathcal{L} .

Prékopa, Ganczer, Deák, Patyi (1980).

Cutting-plane methods approximating the convex level set \mathcal{L} .

Prékopa and Szántai (1978), Szántai (1988), Mayer (1998),

Henrion and associates (2000 -).

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- Efficiency due to reusing former gradient information.
- Difficulty: gradient computation is noisy.
Practicable implementations require sophisticated tolerance handling.

Other solution approaches

Uncertain convex programs,

Campi, Calafiore, Garatti, Caré (2005 -)

Sample Average Approximation, integer programming formulations,

Ahmed, Luedtke, Nemhauser (2007 -)

Cone programming,

Cheng, Gicquel and Lissner (2012).

Classic dual approach

Idea: to build an inner approximation of the level set \mathcal{L} .

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Definition (Prékopa 1990): z is a p -efficient point iff

$$\left\{ \begin{array}{l} z \in \mathcal{L} \quad \text{and} \\ \text{there exists no } z' \in \mathcal{L} \text{ such that } z' \leq z, z' \neq z. \end{array} \right.$$

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Given p -efficient points z_1, \dots, z_k ,

$\text{conv}(z_1, \dots, z_k) + \mathbb{R}_+^n$ is an inner approximation of \mathcal{L} .

Solution methods

Prékopa, Vizvári, Badics (1998):

set of p -efficient points generated before optimization,

Dentcheva, Prékopa, Ruszczyński (2000):

primal-dual method (cone-generation),

Dentcheva, Lai, Ruszczyński (2004), Dentcheva and Martinez (2013).

Cone generation Dentcheva, Prékopa, Ruszczyński (2000)

Primal problem formulated by splitting variables:

$$(\mathcal{P}) \quad \min \mathbf{c}^T \mathbf{x} \quad \text{subject to} \quad T\mathbf{x} = \mathbf{z}, \quad \mathbf{x} \in X, \quad \mathbf{z} \in \mathcal{L}.$$

Lagrangian dual problem obtained by relaxing the constraint $T\mathbf{x} = \mathbf{z}$:

$$(\mathcal{D}) \quad \max_{\mathbf{u}} D(\mathbf{u}) \quad \text{where} \quad D(\mathbf{u}) = \min_{\mathbf{x} \in X} (\mathbf{c}^T - \mathbf{u}^T T)\mathbf{x} \quad + \quad \min_{\mathbf{z} \in \mathcal{L}} \mathbf{u}^T \mathbf{z}.$$

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Solution method:

- from dual viewpoint: cutting-plane method for $D()$,
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New cuts/columns are improving p -efficient points.

These are found by solving subproblems $\min_{\mathbf{z} \in \mathcal{L}} \mathbf{u}^T \mathbf{z}$.

An alternative dual approach: epi-approximation

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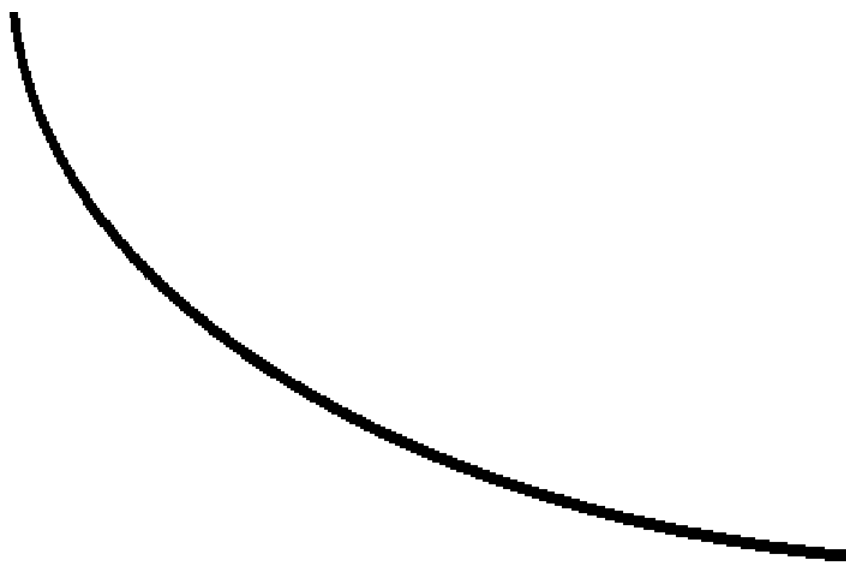
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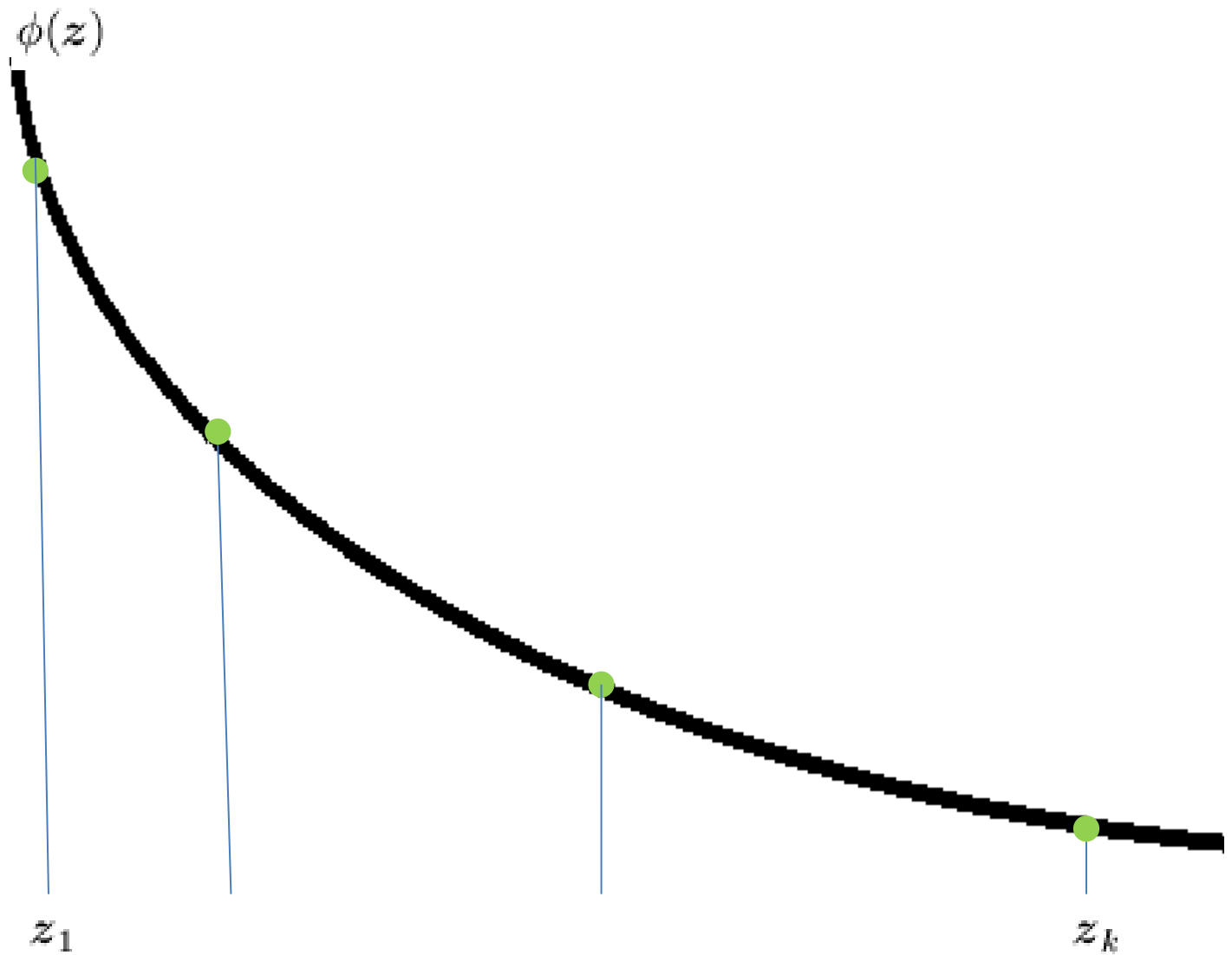
Constraint: $F(\mathbf{z}) \geq p$

Convex formulation: $\underbrace{-\ln F(\mathbf{z})}_{\phi(\mathbf{z})} \leq -\ln p$



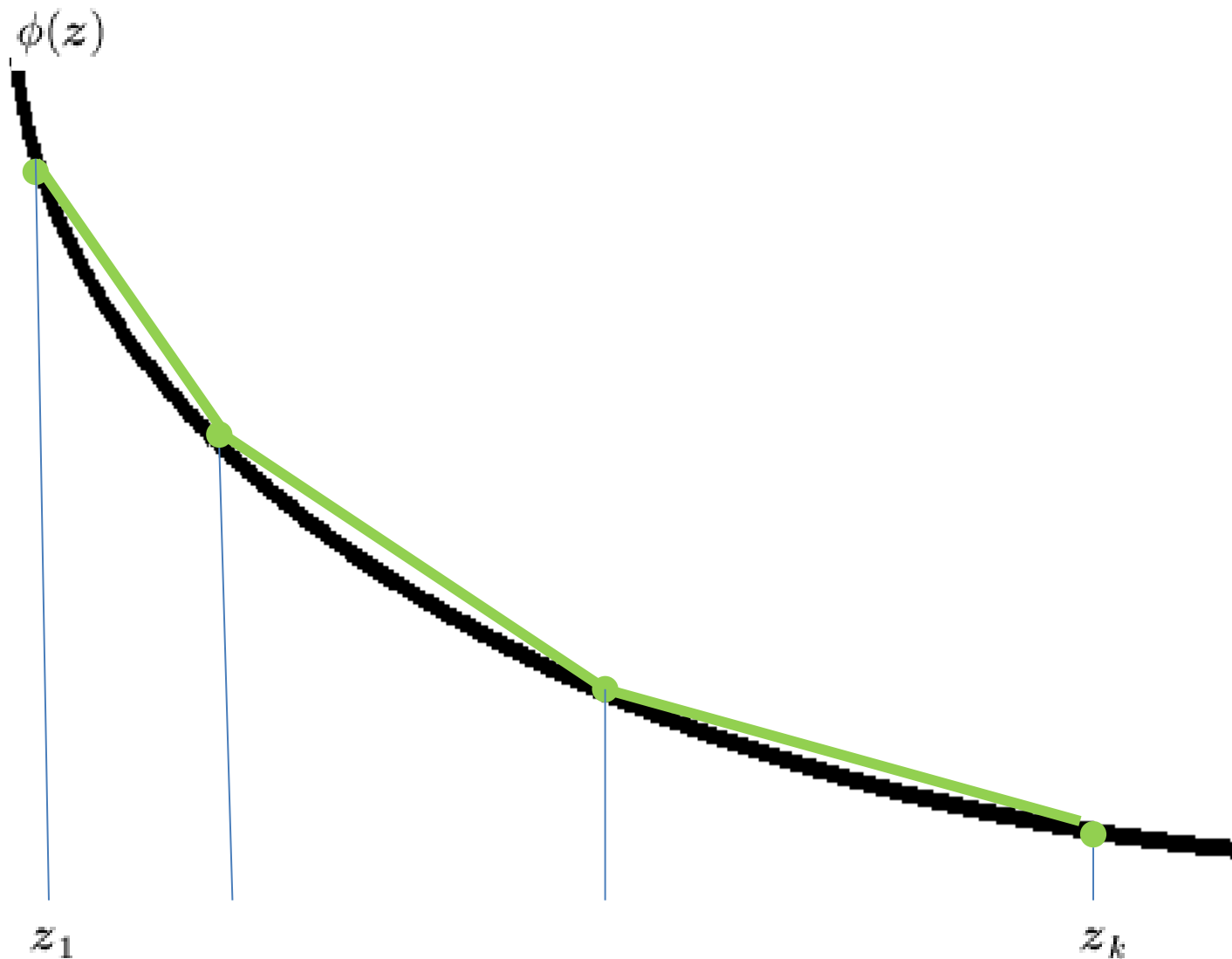
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Epi-approximation to probability maximization

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Epi-approximation computational study

Problems

Normal distribution, dimension up to 15.

Gradients of the distribution function can be computed componentwise:

$$\frac{\partial F(z_1, \dots, z_n)}{\partial z_i} = F(z_1, \dots, z_{i-1}, z_{i+1}, \dots, z_n | z_i) f_i(z_i) \quad (i = 1, \dots, n).$$

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Implementation: column generation scheme

Master problem solved by CPLEX simplex, version 12.6.3.

Subproblems solved by a steepest descent method.

$$\min_{\mathbf{z}} \{ \phi(\mathbf{z}) - \mathbf{u}^T \mathbf{z} \} \quad \text{where} \quad \phi(\mathbf{z}) = -\log F(\mathbf{z}).$$

$F(\mathbf{z})$ function values and gradients computed by Genz's code.

A gradient component can be computed from an appropriate (n-1)-dimensional normal distribution function value.

Experience

Code proved reliable and robust.

Number of iterations depends on problem dimension,
and on the optimal probability achievable.

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Heuristic adjustment

Aim: balancing different efforts:

- solving and resolving master problem (CPLEX),
- solving subproblems (Genz code).

Approximate solution of subproblems.

A single line search made in each steepest descent procedure,
and even this line search is approximate.

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Outcome

Effort spent in a subproblem substantially decreased

Number of master iterations did not increase significantly.

Theoretical justification of heuristic adjustment

An ideal convex programming problem

$$\mathcal{F} = \min_{\mathbf{z}} f(\mathbf{z})$$

Assume the function $f(\mathbf{z})$ is twice continuously differentiable,

and there are $\alpha, \omega \in \mathbb{R}$ ($0 < \alpha \leq \omega$) such that

$$\alpha I \preceq \nabla^2 f(\mathbf{z}) \preceq \omega I \quad (\mathbf{z} \in \mathbb{R}^n).$$

Here the relation $U \preceq V$ means that $V - U$ is positive semidefinite.

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A well-known convergence theorem

We minimize $f(\mathbf{z})$ using a steepest descent method.

Starting from \mathbf{z}^0 , let $\mathbf{z}^1, \dots, \mathbf{z}^j, \dots$ denote the iterates.

Then we have

$$f(\mathbf{z}^j) - \mathcal{F} \leq \left(1 - \frac{\alpha}{\omega}\right)^j [f(\mathbf{z}^0) - \mathcal{F}].$$

convergence theorem

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Application to our column generation scheme

New columns are found by solving subproblems

$$\min_{\mathbf{z}} \{ \phi(\mathbf{z}) - \mathbf{u}^T \mathbf{z} + \text{constant} \}.$$

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$$\mathcal{F} = \min_{\mathbf{z}} \left\{ \underbrace{\phi(\mathbf{z}) - \mathbf{u}^T \mathbf{z}}_{f(\mathbf{z})} + \text{constant} \right\}.$$

Corollary of the theorem

We can find a starting point \mathbf{z}^0 such that

$$\mathbf{u}^T \mathbf{z}^j - \phi(\mathbf{z}^j) - \text{constant} \geq \left[1 - \left(1 - \frac{\alpha}{\omega}\right)^j\right] (-\mathcal{F}).$$

holds with iterate \mathbf{z}^j obtained by j line search iterations.

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Corollary :

We can find a starting point \mathbf{z}^0 such that

$$\underbrace{\mathbf{u}^T \mathbf{z}^j - \phi(\mathbf{z}^j) - \text{constant}}_{\text{reduced cost of column } \mathbf{z}^j} \geq \left[1 - \left(1 - \frac{\alpha}{\omega} \right)^j \right] \underbrace{(-\mathcal{F})}_{\text{largest possible reduced cost value}}$$

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Consider the master problem as a LP problem.

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Consider the master problem as a LP problem.

Even for a moderately large j , the iterate \mathbf{z}^j is

a fairly good improving column in the column generation scheme.

On α, ω

Requirement: $\alpha I \preceq \nabla^2 \phi(\mathbf{z}) \preceq \omega I$

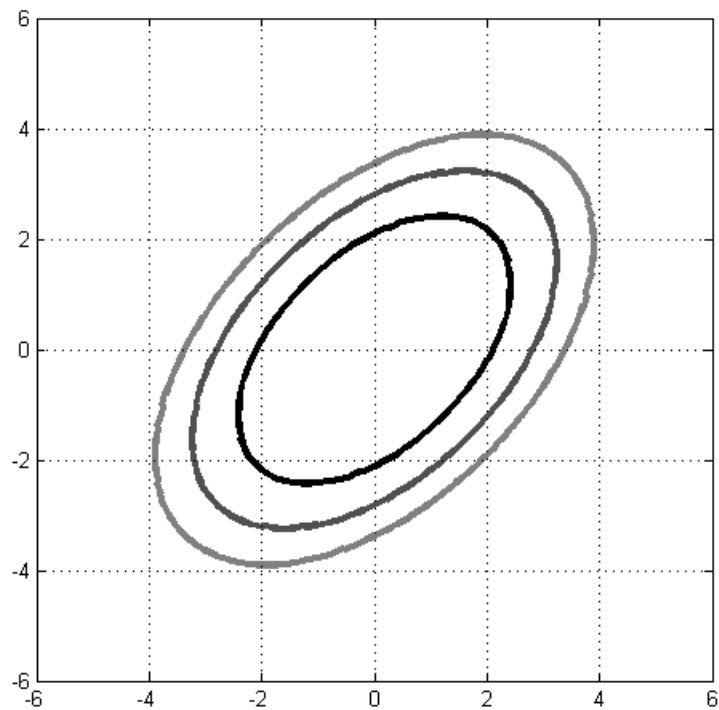
In the present model, we have $\phi(\mathbf{z}) = -\ln F(\mathbf{z})$.

Requirement does not hold for every \mathbf{z} .

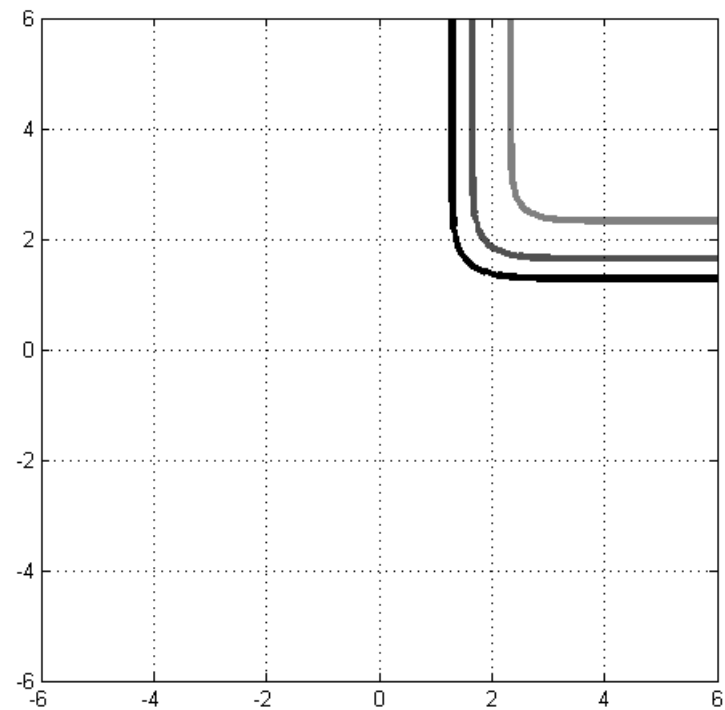
But it holds over a bounded box.

Illustration

Two-dimensional standard normal distribution, covariance = 0.5



Contours of density function



Contours of distribution function

Consider eigenvalues of $-\nabla^2 \log F(\mathbf{z})$.

Larger eigenvalue: $\omega(\mathbf{z})$.

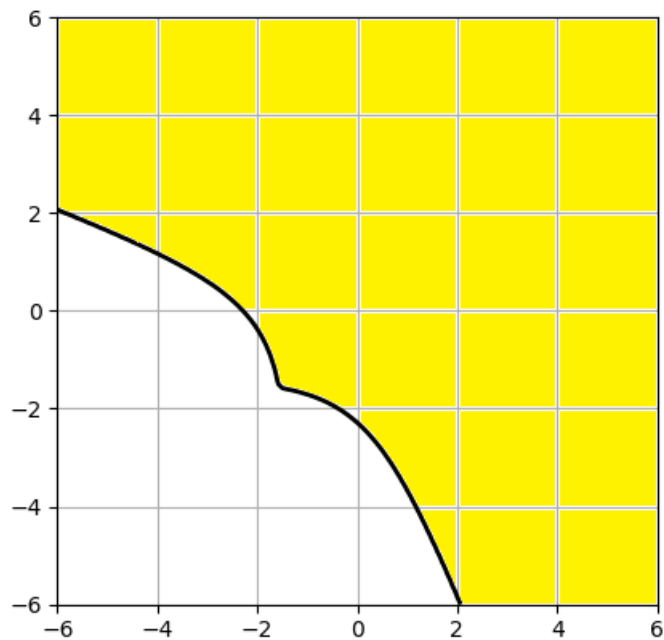
Smaller eigenvalue: $\alpha(\mathbf{z})$.

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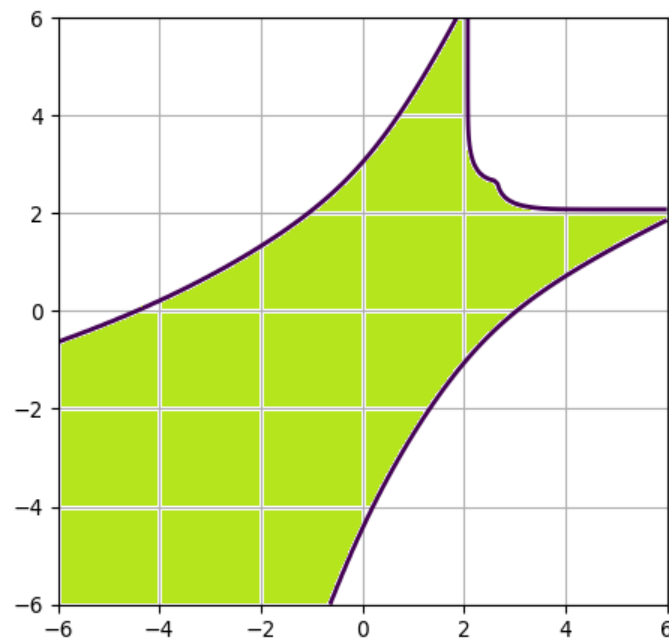
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Level sets of eigenvalue functions



$$\omega(\mathbf{z}) \leq 10$$



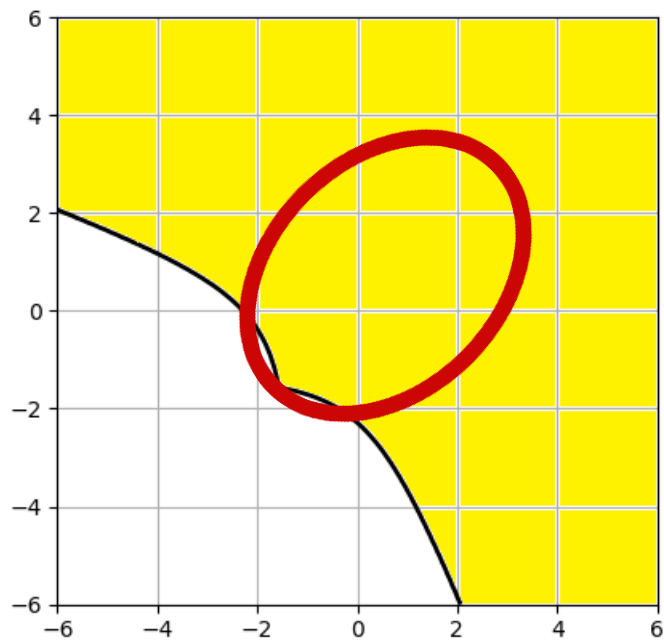
$$\alpha(\mathbf{z}) \geq 0.1$$

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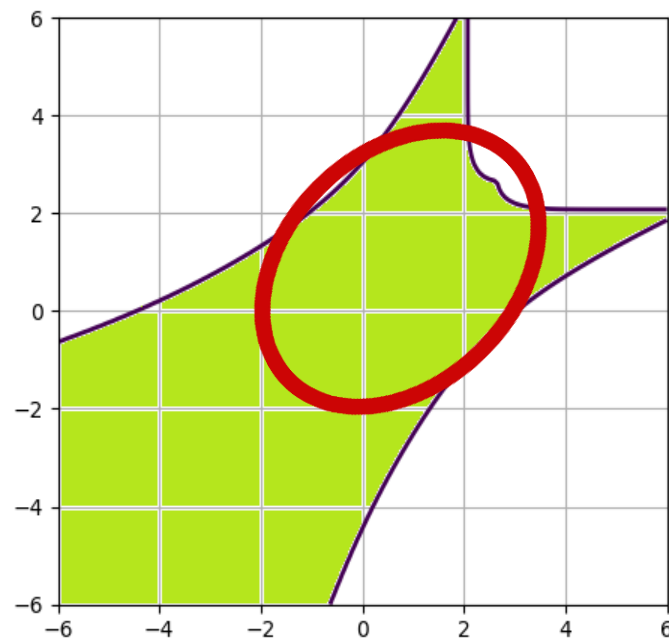
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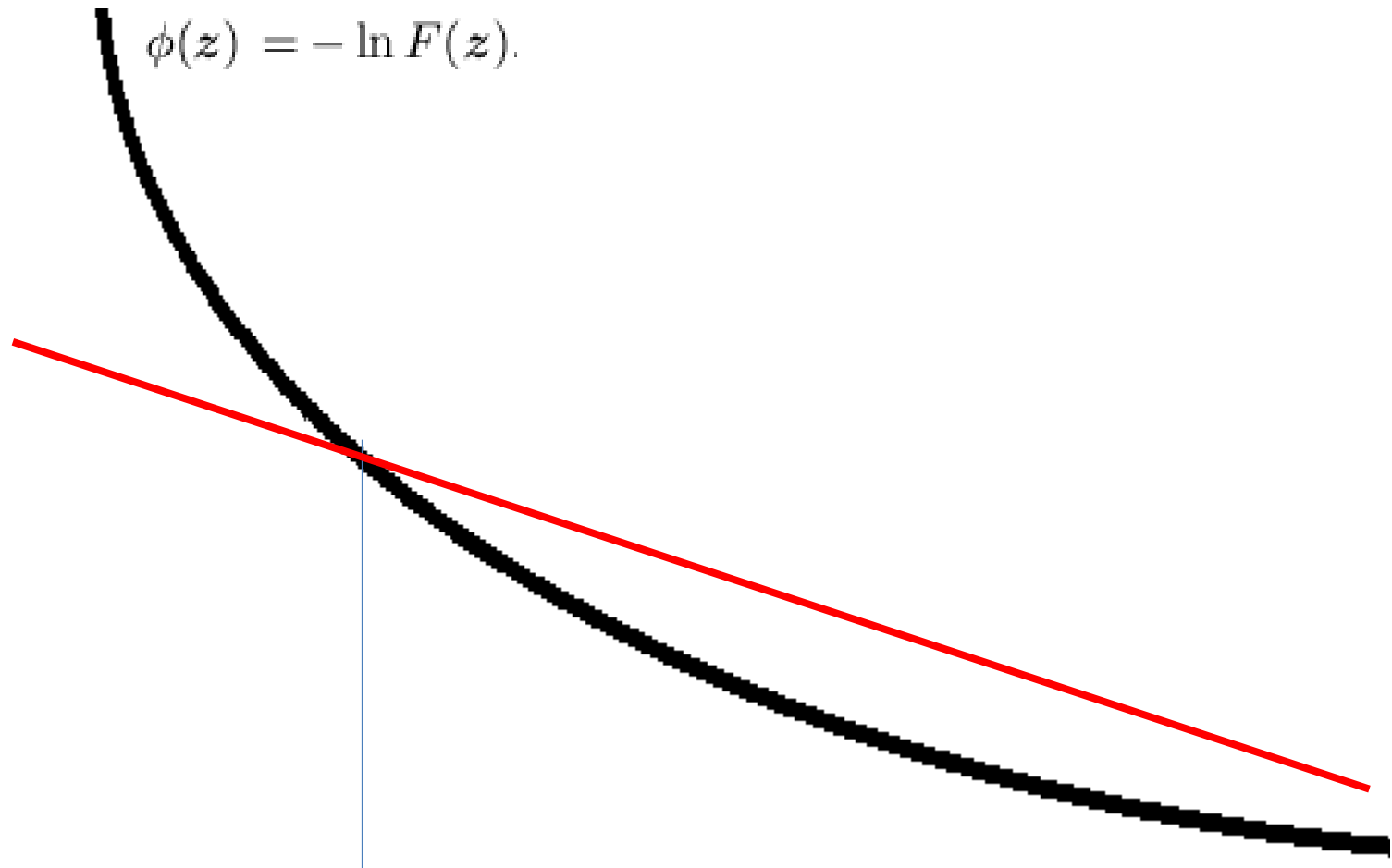
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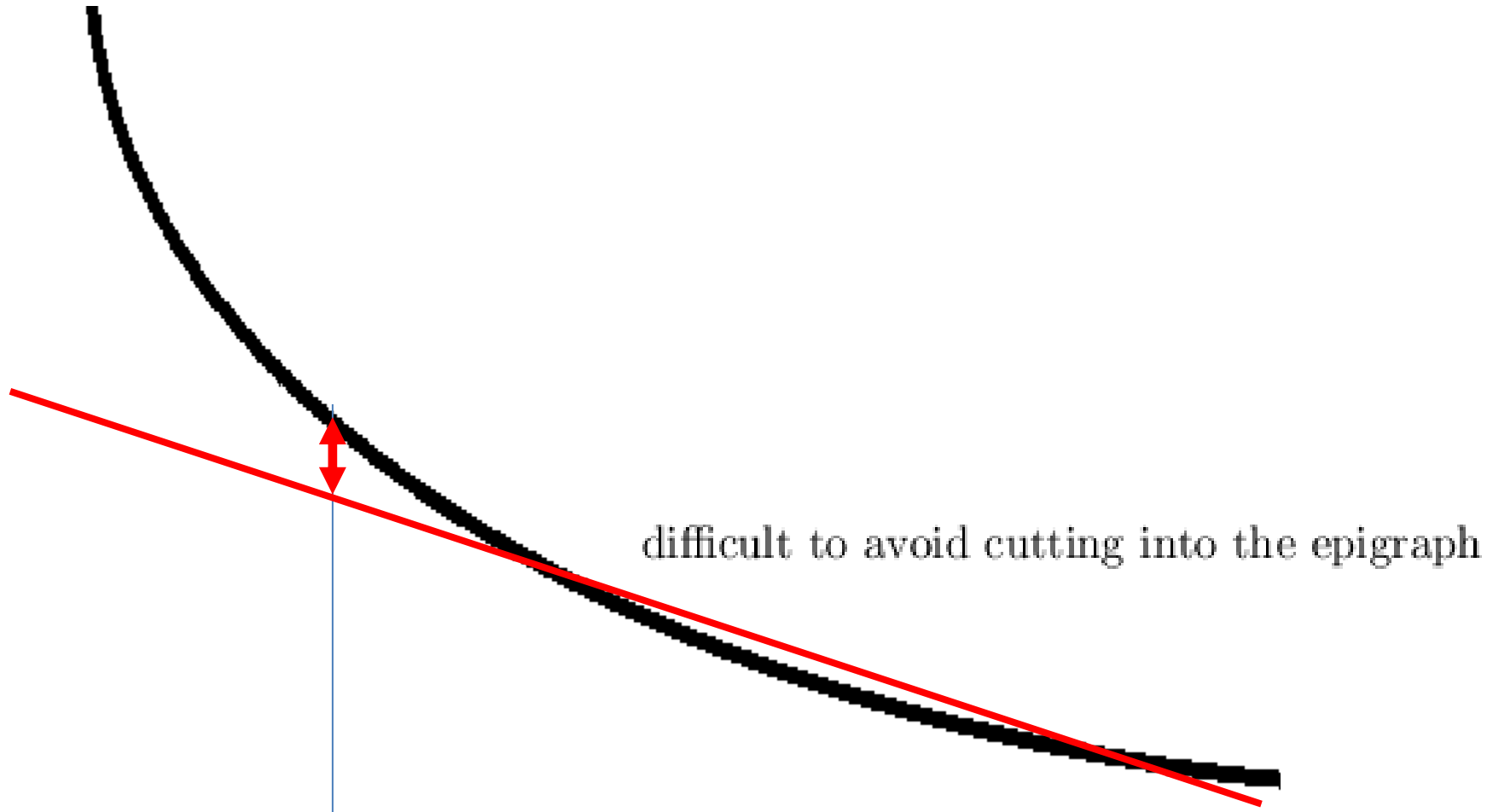
Direct cutting-plane method:



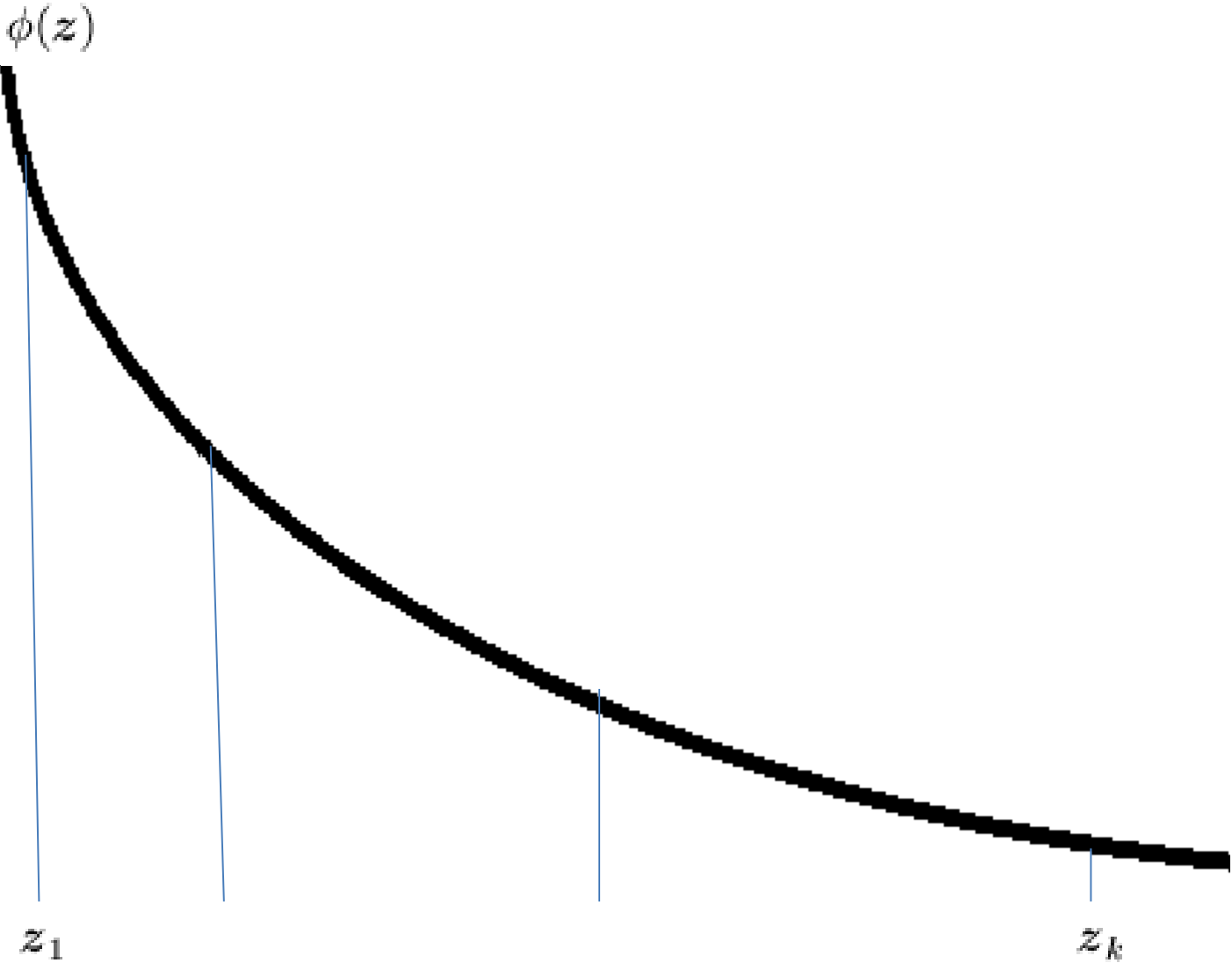
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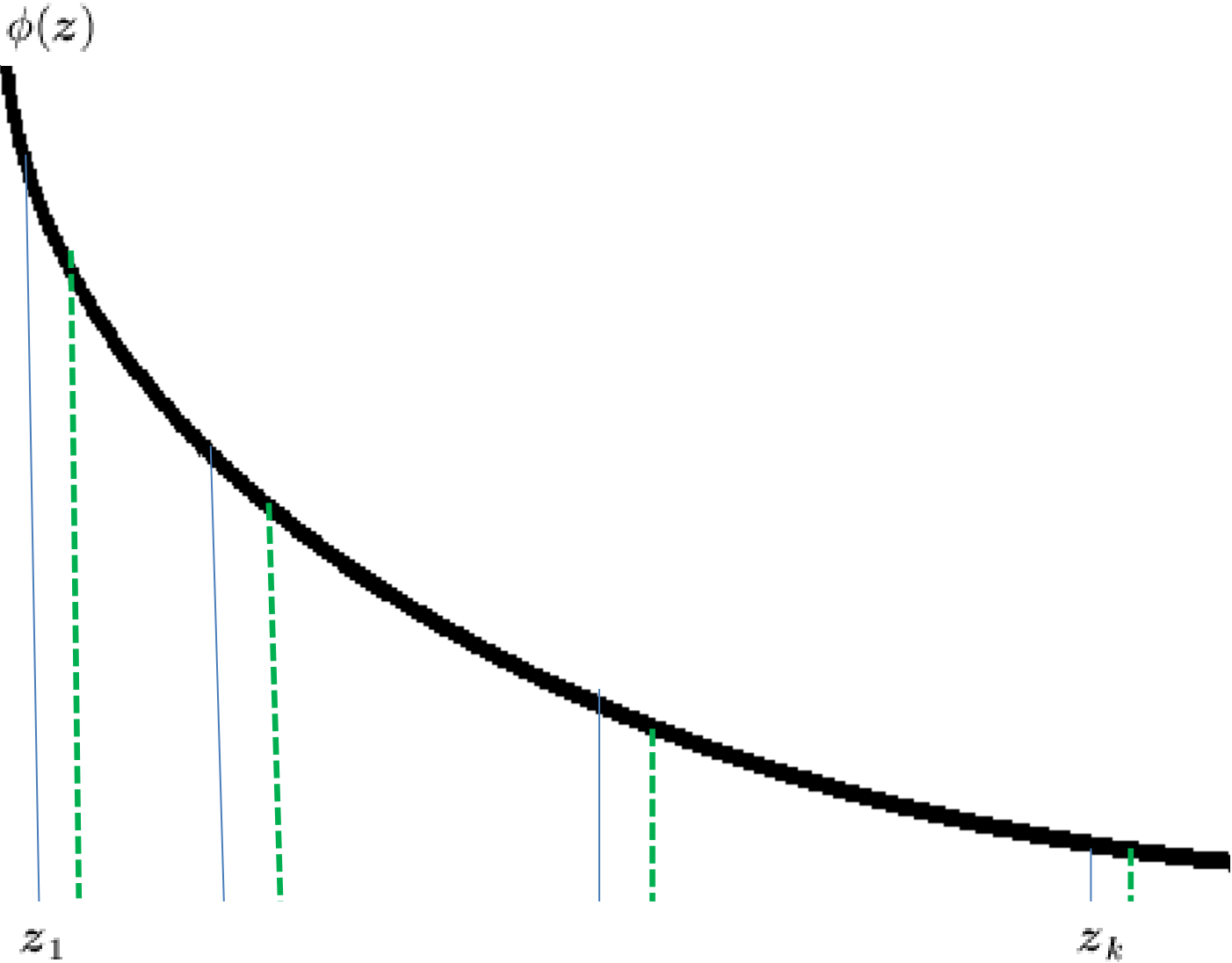
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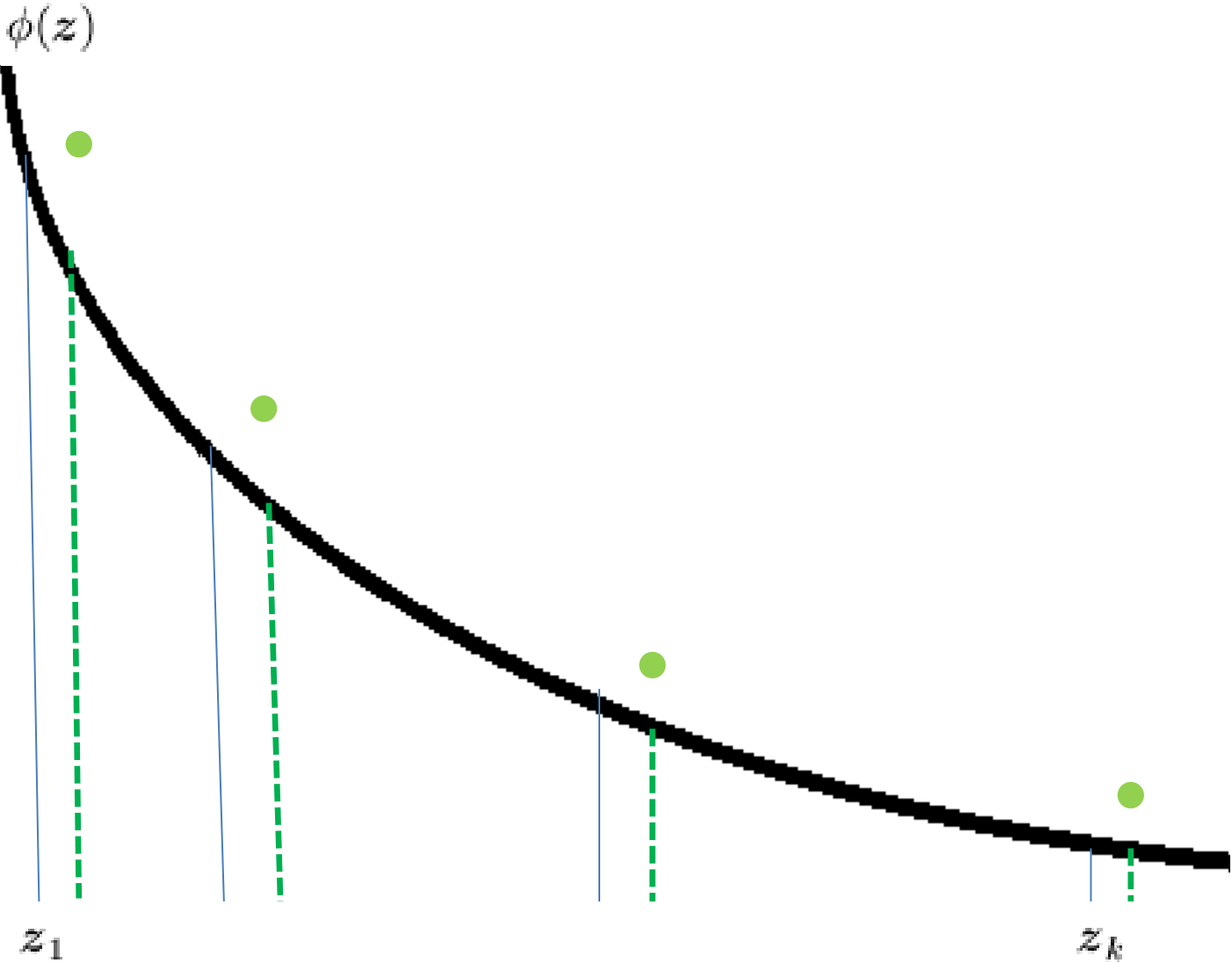
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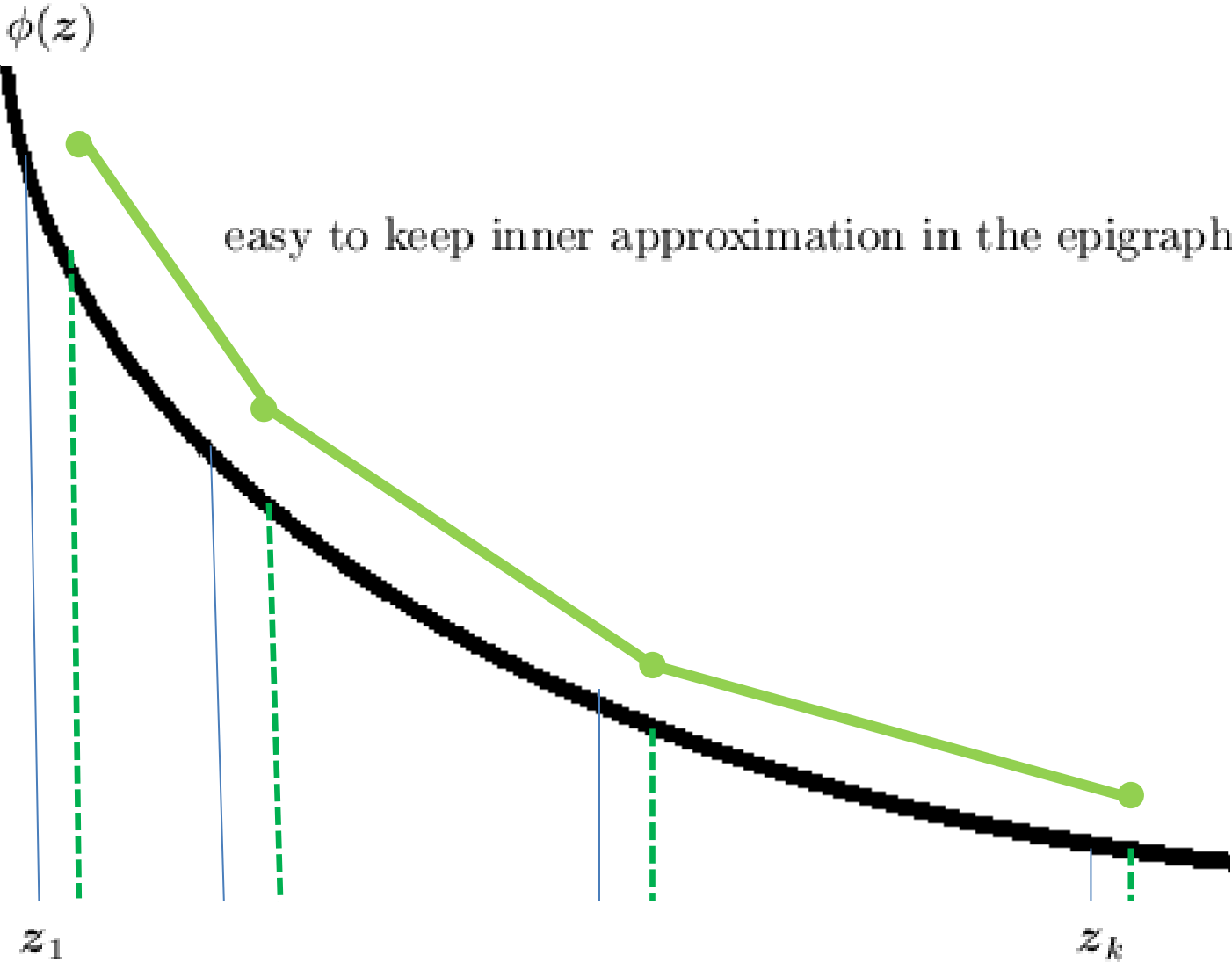
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Working with gradient estimates

Theoretical development for an ideal convex programming problem

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Given $\sigma > 0$, we can construct realizations of a random vector \mathbf{G}° satisfying

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A generalization of the convergence theorem

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Resemblance to the stochastic approximation family

But present approach builds a model problem.

Working with gradient estimates

Application to probability maximization

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Reliable gradient estimates can be constructed using ideas of

Szántai (1976, 1985);

Deák (1980, 1986);

Ambartzumian et al. (1998);

Gassmann (1988); Deák, Gassmann, Szántai (2002);

Mádi-Nagy, Prékopa (2004).

Köszönöm a figyelmet!